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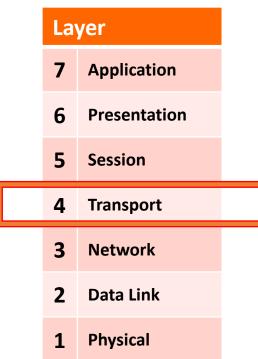
### **Communications and Computer Networks**

Summer Term 2023

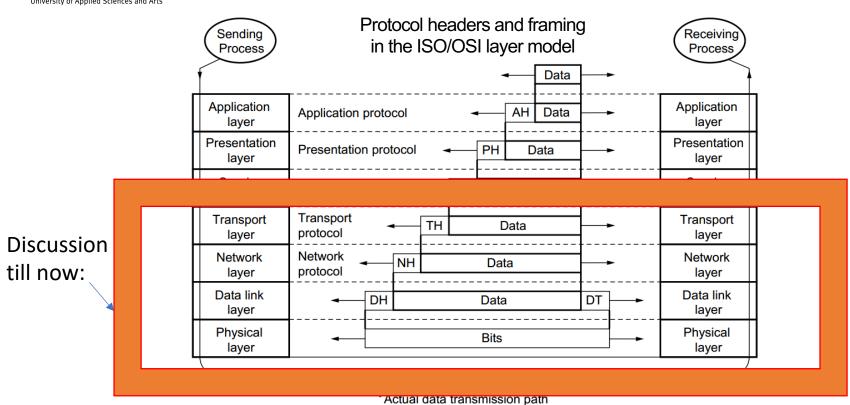
Prof. Dr. Daniel Spiekermann Faculty of Computer Science

### Recap of last lecture (1/8)

- You know the User Datagram Protocol (UDP) and can name its header structure
- You know the Transmission Control Protocol (TCP) and can name its header structure
- You know the principle of and can name selected port
- You know the tasks of TCP and can explain the communication process, including setting up and disconnecting the connection.
- You know the structure of the TCP header and can explain the fields.
- You can explain the principle of TCP window management.
- You are familiar with the Bandwidth-Delay product and can explain its importance for overload control using window control.
- You know the methods fast-retransmit and fast-recovery and can explain them.
- You know the QUIC protocol



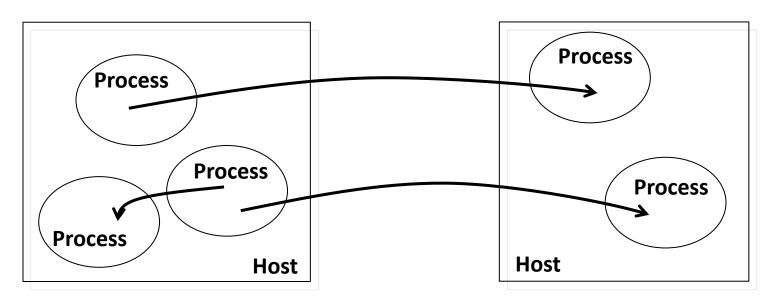
### Recap of last lecture (2/8)





### Recap of last lecture (3/8)

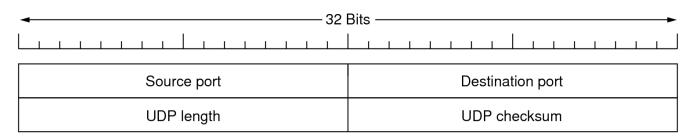
The ultimate goal of the communication mechanism between computers is that a **process** can communicate with another process that runs on another machine (or on the same one) in a transparent way.



## Recap of last lecture (4/8)

Well known service	Reserved Port Number							
TELNET	23							
FTP	20, 21							
DNS	53							
НТТР	80							
HTTPS	443							
SMTP	25							
POP3	110							
IMAP4	143							
IMAP4 over TLS/SSL	993							
SSH	22							
SNMP	161, 162							

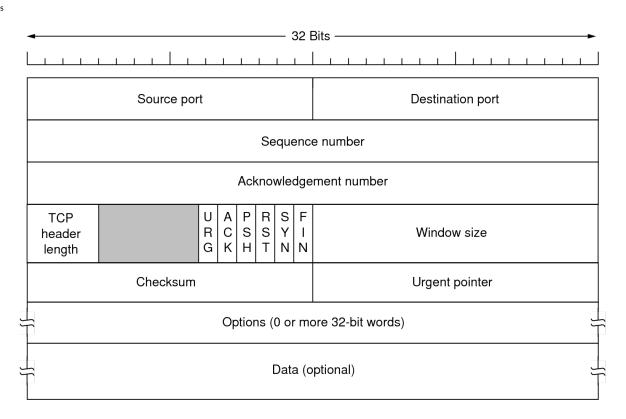
### Recap of last lecture (5/8)



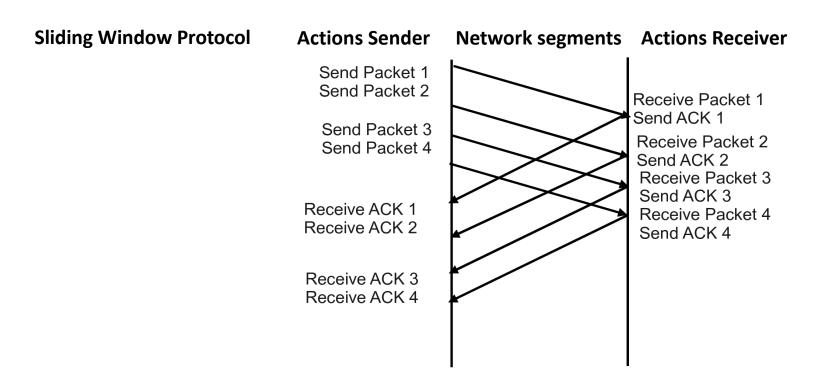
- Source Port: Port number of the sender
- Destination Port: Port number of the recipient
- UDP Length: Length of the datagram including header
- UDP Checksum: Checksum including pseudo-header

The use of Checksum is optional and is used to check the integrity of the header and the packet data. The checksum is calculated by adding a pseudo-header that contains the source and destination address, as well as the protocol number and the UDP packet length, in order to be able to verify without a doubt that the packet reached its correct recipient.

### Recap of last lecture (6/8)



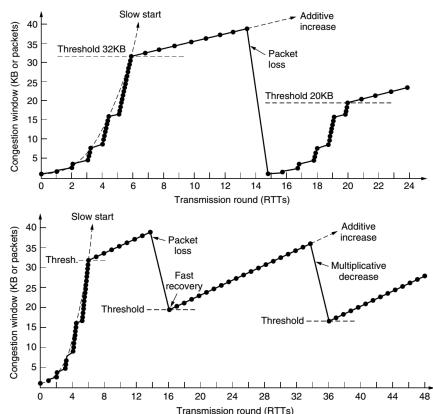
### Recap of last lecture (7/8)



## Recap of last lecture (8/8)

- Slow Start &
- Fast Retransmit

- Slow Start &
- Fast Recovery



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## **Layer 7 - Application Layer**

Prof. Dr. Daniel Spiekermann Faculty of Computer Science

### **Layer 7 – Application layer**

- The last link between application and transmission medium is the application layer.
- **Highest layer** of the OSI-stack
- It is the only layer that has an interface to the application process.
- The application process itself is "outside" the scope of the layer model.
- Relevant protocols:
  - o DNS
  - o DHCP
  - o HTTP
  - o SMTP
  - 0 ...

#### **Content Overview**

- You know the task of DHCP
- You can describe the process and the involved DHCP messages
- You know the DHCPv6 protocol
- You understand the difference between stateful and stateless addresses.
- You know DNS and the process of name resolution
- You understand the hierarchy of names, TLDs and SLDs as well as FQDNs
- You know HTTP and HTTPS
- You understand the request response communication and the interaction of application and protocol

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### **Application Layer - DHCP**

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### Static ip address assignment

```
$ ip address add 192.168.11.1/24 dev eth0
$ ifconfig eth0 192.168.11.1 netmask 255.255.255.0
$ cat /etc/network/interfaces
iface eth0 inet static
address 192.168.11.1
netmask 255,255,255.0
gateway 192.168.11.254
```

### **Dynamic Host Configuration Protocol (RFC 2131)**

- DHCP (Dynamic Host Configuration Protocol) enables the dynamic assignment of an IP address (IPv4) and other configuration parameters to computers on a network using a DHCP server.
- Using DHCP, a server can automatically assign parameters such as IP address, netmask, gate-way and DNS server when starting a new computer
- DHCP is based on the BOOTP protocol, which can be used to implement driveless workstations
  that first get an IP address from the BOOTP server, and then reload a bootable operating system
  from the network, with which they then boot.
- The technology eliminates the need for individually configuring network devices manually, and consists of two network components, a centrally installed network DHCP server and client instances of the protocol stack on each computer or device. When connected to the network, and periodically thereafter, a client requests a set of parameters from the DHCP server using the DHCP protocol.
- DHCP services exist for networks running Internet Protocol version 4 (IPv4), as well as version 6 (IPv6). The IPv6 version of the DHCP protocol is commonly called DHCPv6.

### **DHCP Operation**

- The DHCP employs a **connectionless** service model, using the User Datagram Protocol (**UDP**). It is implemented with two UDP port numbers for its operations which are the same as for the bootstrap protocol (BOOTP). UDP port number **67** is the destination port of a **server**, and **UDP** port number **68** is used by the client.
- DHCP operations fall into four phases: server discovery, IP lease offer, IP lease request, and IP lease acknowledgement. These stages are often abbreviated as DORA for discovery, offer, request, and acknowledgement.
- The DHCP operation begins with clients broadcasting a request. If the client and server are in different Broadcast Domains, a DHCP Helper or **DHCP Relay Agent** may be used. Clients requesting renewal of an existing lease may communicate directly via UDP unicast, since the client already has an established IP address at that point.

#### **DHCP-Server allocation**

- Listens on port 67 UDP
- Provides at least a range of IP-addresses

#### Dynamic allocation

A network administrator reserves a range of IP addresses for DHCP, and each DHCP client on the LAN is configured to request an IP address from the DHCP server during network initialization.

```
subnet 192.168.1.0 netmask 255.255.255.0 {
range 192.168.1.10 192.168.1.100;
range 192.168.1.150 192.168.1.200;
}
```

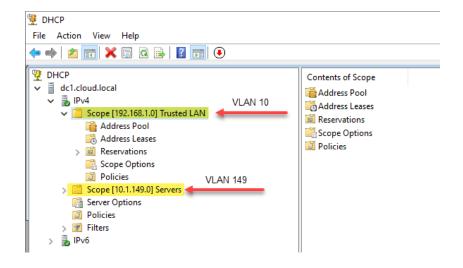
#### Manual allocation

This method is also variously called static DHCP allocation, fixed address allocation, reservation, and MAC/IP address binding. An administrator maps a unique identifier (a client id or MAC address) for each client to an IP address, which is offered to the requesting client.

```
# always set IP to 10.0.3.1 for web1 vm
10.0.3.1,web1
# always set IP to 10.0.5.12 for db12 vm based on mac
00:16:ee:62:ff:90,10.0.5.12,db12
```

#### **DHCP Server**

- Typically every DSL-Router provides a DHCP-server
- Dedicated DHCP-server for OS exist
  - Windows Server
  - dnsmasq
  - isc-dhcpd (deprecated since 2022)
  - o Kea



```
# Sample configuration file for ISC dhcpd for Debian

# Attention: If /etc/ltsp/dhcpd.conf exists, that will be used as

# configuration file instead of this file.

authoritative;
default-lease-time 180;
max-lease-time 7200;
option routers 192.168.1.10;
option domain-name-servers 192.168.1.1, 192.168.1.2;
option domain-name "EMEA.com";

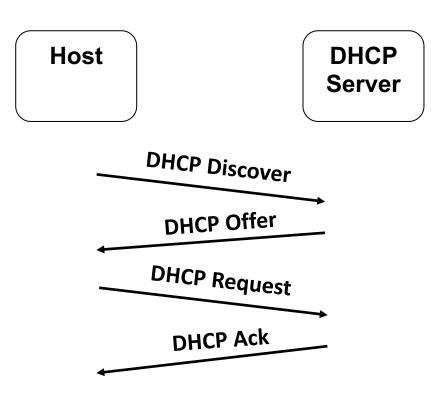
subnet 192.168.1.0 netmask 255.255.255.0 {
   range 192.168.1.10 192.168.1.200;
}

AG Get Help AD WriteOut AR Read File AY Prev Page AK Cut Text AC Cur Pos

XX Exit AJ Justify Where Is AV Next Page AU Uncut TextAT To Spell
```

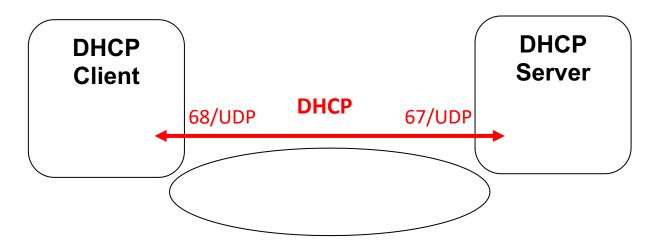
#### **DHCP DORA**

- DHCP Discover broadcast to locate servers
- DHCP Offer
   each server offers configuration
   parameters
- DHCP Request request offered parameters from one server
- DHCP Ack assignment confirmation



#### **DHCP and Wireshark**

Source	Destination	Protocol Length	Info
0.0.0.0	255.255.255.255	DHCP 618	DHCP Discover
192.168.0.1	255.255.255.255	DHCP 342	DHCP Offer
0.0.0.0	255.255.255.255	DHCP 618	DHCP Request
192.168.0.1	255.255.255.255	DHCP 342	DHCP ACK



#### **DHCP** structure

0	0	2	0 3	0 4	0 5	0	0 7	0	0 9	1 0	1 1	1 2	1 3	1 4	1 5	1 6	1 7	1 8	1 9	2 0	2	2 2	2 3	2 4	2 5	2 6	2 7	2 8	2 9	3	3
	OP HTYPE												HLEN HOPS																		
XID											ID .																				
SECS													FLAGS																		
														(	CΙΑ	DD	R														
														`	ΛIΛ	DD	R														
														5	SIA	DD	R														
	GIADDR																														
	CHADDR																														
	SNAME																														
	FILE																														
	OPTIONS																														

#### Wireshark and DHCP

```
Ethernet II, Src: cc:01:0a:c4:00:00 (cc:01:0a:c4:00:00), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
Internet Protocol Version 4, Src: 192.168.0.1, Dst: 255.255.255.255
> User Datagram Protocol, Src Port: 67, Dst Port: 68
Dynamic Host Configuration Protocol (Offer)
   Message type: Boot Reply (2)
    Hardware type: Ethernet (0x01)
   Hardware address length: 6
   Hops: 0
    Transaction ID: 0x0000155c
    Seconds elapsed: 0
   Bootp flags: 0x8000, Broadcast flag (Broadcast)
    Client IP address: 0.0.0.0
    Your (client) IP address: 192.168.0.3
   Next server IP address: 0.0.0.0
    Relay agent IP address: 0.0.0.0
    Client MAC address: cc:00:0a:c4:00:00 (cc:00:0a:c4:00:00)
```

#### Wireshark and DHCP

**Broadcast** 

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Unique MAC from sender

```
Ethernet II, Src: cc:01:0a:c4:00:00 (cc:01:0a:c4:00:00), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
> Internet Protocol Version 4, Src: 192.168.0.1, Dst: 255.255.255.255
 User Datagram Protocol, Src Port: 67, Dst Port: 68
Dynamic Host Configuration Protocol (Offer)
    Message type: Boot Reply (2)
    Hardware type: Ethernet (0x01)
   Hardware address length: 6
   Hops: 0
                                                                        Port 67 und 68
    Transaction ID: 0x0000155c
    Seconds elapsed: 0
   Bootp flags: 0x8000, Broadcast flag (Broadcast)
    Client IP address: 0.0.0.0
    Your (client) IP address: 192.168.0.3
   Next server IP address: 0.0.0.0
    Relay agent IP address: 0.0.0.0
    Client MAC address: cc:00:0a:c4:00:00 (cc:00:0a:c4:00:00)
```

#### **DHCP structure**

- OP: opcode 1 = REQUEST, 2 = REPLY
- HTYPE: Hardware address type (01 p/MAC Addresses)
- HLEN: Hardware address length (06 p/MAC Addresses)
- HOPS: Client sets to zero, optionally used by relay agents
- XID: Transaction ID
- SECS: seconds elapsed since client began addr acquisition
- FLAGS: BROADCAST flag
- CIADDR: Client IP address
- YIADDR: 'your' (client) IP address, the offer of the server
- **SIADDR**: IP address of server
- **GIADDR**: Relay agent IP address
- **CHADDR**: Client hardware address
- SNAME: Optional server host name

### **DHCP Parameters and Options**

- DHCP provides assignment of further parameters (so-called options)
  - Subnet Mask (1)
  - Default Gateway (3)
  - Time-Server (4)
  - DNS Server (6)
  - Hostname (12)
  - End (255)

```
> Option: (12) Host Name
> Option: (1) Subnet Mask (255.255.255.0)

Option: (3) Router
   Length: 4
   Router: 192.168.0.1

Option: (6) Domain Name Server
   Length: 8
   Domain Name Server: 192.168.0.1
   Domain Name Server: 192.168.1.1

> Option: (255) End
```

#### **DHCP** Lease

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- Different times in DHCP ACK
- LEASE time (Option 51)
  - Time period for which the assignment is valid.
- Renewal Time (50% of the length of the lease)
  - This option specifies the time interval from address assignment until the client transitions to the RENEWING state. Client sends DHCPREQUEST
- Rebinding Time (87.5% (7/8ths) of the length of the lease)
  - This option specifies the time interval from address assignment until the client transitions to the REBINDING state.
     When the server does not reply to the DHCPREQUEST, client tries to connect to any DHCP-server in the network

Option: (51) IP Address Lease Time
Length: 4

IP Address Lease Time: (60s) 1 minute

Option: (58) Renewal Time Value
Length: 4

Renewal Time Value: (30s) 30 seconds

∨ Option: (59) Rebinding Time Value

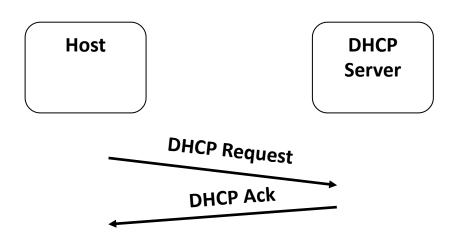
Length: 4

Rebinding Time Value: (52s) 52 seconds

### **DHCP Request**

#### **DHCP Request**

- Send to the server after renewal timer reached 0
- Send directly to the server, no broadcast
- Server answers with DHCP ACK, if requested IP-address is still available for the host



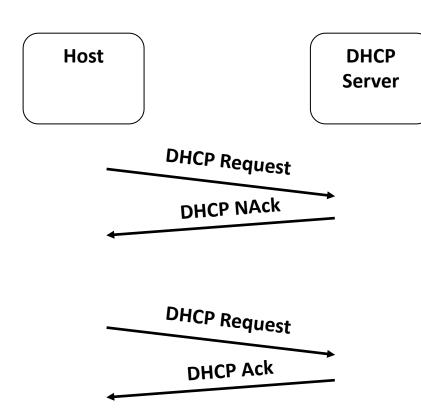
#### **DHCP NACK**

DHCP NACK

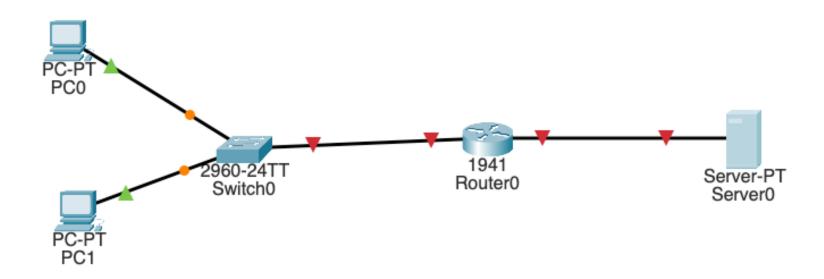
Typical case of network change.

The host already has an IP and when reconnecting to a network, it requests the renewal of the Lease with a DHCP Request, but the IP is no longer valid in that network.

Host sends a new DHCP Request to the server

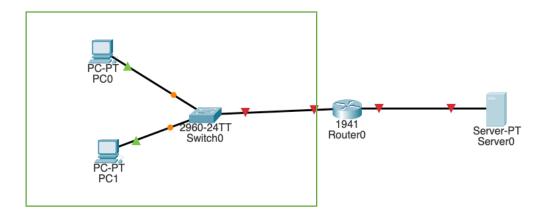


## A "new" problem



### **DHCP Relaying**

- The DHCP client broadcasts on the local link, the client's broadcast can only be received on its own subnet.
- DHCP relay agents can be installed on these subnets. A DHCP relay agent runs on a network device, capable of routing between the client's subnet and the subnet of the DHCP server.



### **DHCP Relaying**

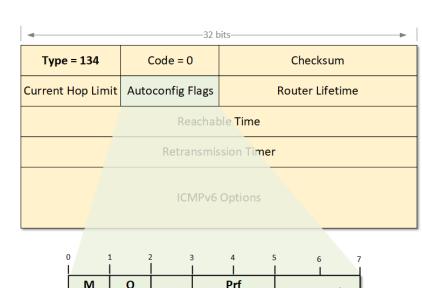
- The DHCP client broadcasts on the local link, the client's broadcast can only be received on its own subnet.
- **DHCP relay agents** can be installed on these subnets. A DHCP relay agent runs on a network device, capable of routing between the client's subnet and the subnet of the DHCP server.
- The DHCP server uses the GIADDR-value to determine the subnet, and subsequently the corresponding address pool, from which to allocate an IP address. When the DHCP server replies to the client, it sends the reply to the GIADDR-address, again using unicast.
- The relay agent then retransmits the response on the local network, using unicast (in most cases) to the newly reserved IP address, in an ethernet frame directed to the client's MAC address.
- The communication between the relay agent and the DHCP server typically uses both a source and destination **UDP** port of **67**.

### **DHCPv6 (RFC 3315, now RFC 8415)**

- Uses UDP port 546 (client) and 547 (server)
- DHCP server (and relay agents) listen on ff02::1:2 (Multicast) or ff05::1:3
- Client requests
- Stateless:
  - Prefix is assigned by RA
  - o DHCPV6 provides network information like DNS server
- Stateful:
  - RA with flag "managed"
  - Full ip address configuration assigned by the server
- Server and client have a DUID (DHCP Unique Identifier)
  - DHCP servers use DUIDs to identify clients for the selection of configuration parameters and in the association of IAs with clients.
  - o DHCP clients use DUIDs to identify a server in messages where a server needs to be identified.
- LLA is still calculated by the client itself

#### DHCPv6 stateful with RA

- Needs specific configuration of the router
   set M-flag = 1
- ICMPv6 Type134 (RA) is send after RS
- Client is informed to get all information from a DHCPv6 server



flag

flag

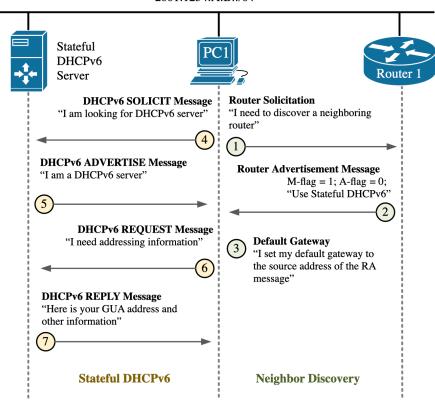
flag

Reserved

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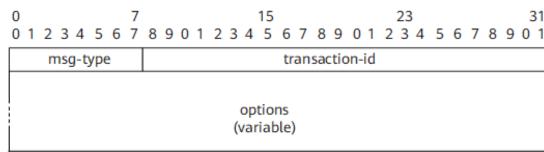
#### DHCPv6 stateful with RA

2001:1234:A:B::/64



#### **DHCPv6 Protocol format**

- Msg-type: Message Type
- Transaction-id: Identifies transaction between client and server
- Options: Indicates an option field
  - Option-code
    - 1: Client-ID
    - 2: Server-ID
    - 3: Identity Association for Non-temporary Addresses
    - 18: IID
  - Option-length
  - Option-data





### **DHCPv6 Message Types**

#### 13 different msg-types

Туре	Packet	DHCPv4 mapping
1	Solicit	Discover
2	Advertise	Offer
3	Request	Request
4	Confirm	-
5	Renew	Request
6	Rebind	Request
7	Reply	ACK/NACK

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#### **DHCPv6** and Wireshark

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```
1.) Request to all servers or relay agents
                                                                                  2.) Answer of a server intended, but LLA
                                                                                  not known, so NS send
fe80::a00:27ff:fefe:8f95
                              ff02::1:2
                                                             DHCPV6
                                                                           114 UDP
                                                                                          Solicit XID: 0x100874 CID: 000100011c39cf8
fe80::a00:27ff:fed4:10bb
                              ff02::1:fffe:8f95
                                                             ICMPv6
                                                                            86 ICMPv6
                                                                                          Neighbor Solicitation for fe80::a00:27ff:
fe80::a00:27ff:fefe:8f95
                              fe80::a00:27ff:fed4:10bb
                                                             ICMPv6
                                                                            86 ICMPv6
                                                                                          Neighbor Advertisement fe80::a00:27ff:fefc
fe80::a00:27ff:fed4:10bb
                              fe80::a00:27ff:fefe:8f95 \
                                                             DHCPv6
                                                                           147
                                                                               UDP
                                                                                          Advertise XID: 0x100874 CID: 000100011c39
fe80::a00:27ff:fed4:10bb
                              ff02::16
                                                             ICMPv6
                                                                               IPv6 Hop...
                                                                                          Multicast Listener Report Message v2
fe80::a00:27ff:fefe:8f95
                              ff02::1:2
                                                             DHCPv6
                                                                               UDP
                                                                                           Reguest XID: 0x49174e CID: 000100011c39cf8
                                                                           161
fe80::a00:27ff:fed4:10bb
                              fe80::a00:27ff:fefe:8f95
                                                             DHCPv6
                                                                           147 UDP
                                                                                           Reply XID: 0x49174e CID: 000100011c39cf880
```

#### 3.) Client answers with NA with its LLA

LLA wit

Ethernet II, Src: PcsCompu\_d4:10:bb (08:00:27:d4:10:bb), Dst: PcsCompu\_fe:8f:95 (08:00:27:fe:
Internet Protocol Version 6. Src: fe80::a00:27ff:fed4:10bb. Dst: fe80::a00:27ff:fefe:8f95

```
Internet Protocol Version 6, Src: fe80::a00:27ff:fed4:10bb, Dst: fe80::a00:27ff:fefe:8f95
User Datagram Protocol, Src Port: 547, Dst Port: 546
DHCPv6
  Message type: Advertise (2)
  Transaction ID: 0x100874
Identity Association for Prefix Delegation
    Option: Identity Association for Prefix Delegation (25)
    Lenath: 41
     IAID: 27fe8f95
    T1: 0
    T2: 0
  V IA Prefix
       Option: IA Prefix (26)
       Length: 25
       Preferred lifetime: 4500
       Valid lifetime: 7200
       Prefix length: 64
       Prefix address: 2001:0:0:fe00::
```

4.) Server sends DHCPv6 on port 546 to LLA with the requested information

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## Fachhochschule Dortmund

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## **Application Layer - DNS**

Prof. Dr. Daniel Spiekermann Faculty of Computer Science



# The Domain Name System (RFCs 1034, 1035, 1591)

You want to communicate with www.fh-dortmund.de

What do we need?



# The Domain Name System (RFCs 1034, 1035, 1591)

You want to communicate with www.fh-dortmund.de

We have learned, that communication "uses" ip-addresses



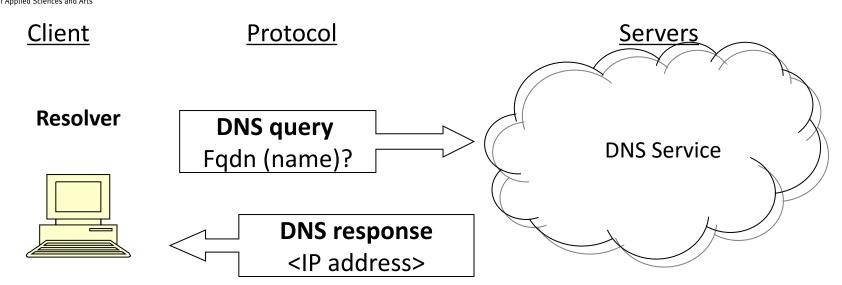
# The Domain Name System (RFCs 1034, 1035, 1591)

You want to communicate with www.fh-dortmund.de

We have learned, that communication "uses" ip-addresses

## So, what is the current ipv4 or ipv6 address of www.fh-dortmund.de?

## **DNS lookups**



## **Protocol format**

ID: 16-bit value that is used to identify the DNS message

- QR (query/response):
  - 0 = request, 1 = response type.
- **opcode:** type of query carried by a message. This field value is repeated in the response.
- AA: Authoritative Answer.
- TC: Truncation
- RD: Recursion Desired
- RA: Recursion Available
- **Zero:** reserved, set to 0.
- **rCode:** Response Code. query was answered successfully or not.

# 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 QR Opcode AA TC RD RA Z RCODE QDCOUNT ANCOUNT ARCOUNT ARCOUNT

**Header Format** 

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## Wireshark - DNS Request

University of Applied Sciences and Arts User Datagram Protocol, Src Port: 55719, Dst Port: 53 Domain Name System (query) Transaction ID: 0x7cf4 Flags: 0x0100 Standard query 0... = Response: Message is a query .000 0... .... = Opcode: Standard query (0) .... ..0. .... = Truncated: Message is not truncated .... 1 .... = Recursion desired: Do query recursively .... = Z: reserved (0) .... = Non-authenticated data: Unacceptable Ouestions: 1 Answer RRs: 0 Authority RRs: 0 Additional RRs: 0 v Oueries www.neverssl.com: type A, class IN Name: www.neverssl.com [Name Length: 16] [Label Count: 3]

Type: A (Host Address) (1)

Class: IN (0x0001)

[Response In: 116]

44

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[Request In: 115]

[Time: 0.256504000 seconds]

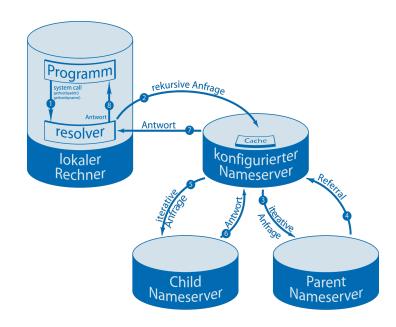
## Wireshark – DNS response

```
University of Applied Domain Name System (response)
           Transaction ID: 0x7cf4
         Flags: 0x8180 Standard query response, No error
             1... = Response: Message is a response
             .000 0... = 0pcode: Standard guery (0)
             .... 0.. .... = Authoritative: Server is not an authority for domain
             .... ..0. .... = Truncated: Message is not truncated
             .... = Recursion desired: Do guery recursively
             .... 1... = Recursion available: Server can do recursive gueries
             .... = Z: reserved (0)
             .... .... .0. .... = Answer authenticated: Answer/authority portion was not a
             .... = Non-authenticated data: Unacceptable
             .... .... 0000 = Reply code: No error (0)
           Ouestions: 1
           Answer RRs: 1
           Authority RRs: 0
           Additional RRs: 0
          > Oueries
         Answers
           www.neverssl.com: type A, class IN, addr 34.223.124.45
               Name: www.neverssl.com
               Type: A (Host Address) (1)
               Class: IN (0x0001)
               Time to live: 49 (49 seconds)
               Data length: 4
               Address: 34.223.124.45
```

/orks

## Types of DNS requests and responses

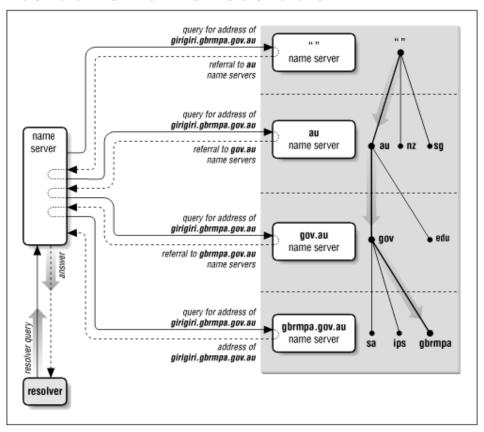
- Authoritative: A name server responds with the information from its local zone file.
- Recursive: The name server itself sends requests to other nameservers before responding. (The requester does not see these requests.)
- Iterative: The requested name server does not send a direct response but points to another name server (forwarding). Root servers answer requests only iteratively.



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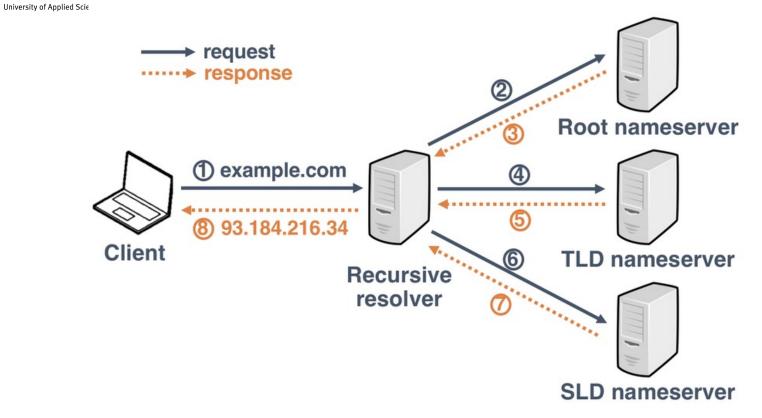
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## **Iterative name resolution**



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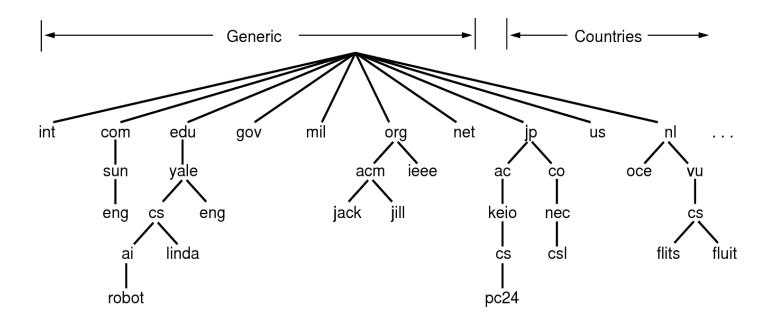
## **Recursive**



#### **DNS** overview

- It is a single distributed database
- The database is **distributed** among many individual servers, which form a federation of nodes that collectively provide the DNS service (the DNS service is implemented in a distributed, not centralized manner)
- Registrations, cancellations and modifications to the data are also done decentralized. The data set is divided into parts called zones. For each zone, an administrative entity is designated that has permission (authority) to modify data contained in the zone.
- The DNS data is structured according to a hierarchical model (tree) to organize the data

## **DNS** hierarchy



#### **DNS** services

## Name resolution

- O Name -> IP-address
- owww.fh-dortmund.de -> 193.25.16.26
- oand vice versa

## Load Balancing

- A name can map to multiple hosts thus multiple addresses
- DNS server returns all addresses but rotates ordering
- Mail server addressing
  - Special mail server addressing apart from normal name resolution

## **DNS Resource Records**

Field	Description	Length (octets)
NAME	fully qualified domain name of the node to which this record pertains	Variable
ТҮРЕ	Type of RR in numeric form (e.g., 1 for A)	2
CLASS	Class code (set to IN for Internet)	2
TTL	Time To Live Seconds that the RR could stays valid in cache	4
RDLENGTH	Length of RDATA field (specified in octets)	2
RDATA	Additional RR-specific data	Variable

HTTPS

## **DNS RR Type**

Type: 65

■ A	Type: 1	RData: IPv4 address
■ NS	Type: 2	RData: FQDN Name Server
<ul><li>CNAME</li></ul>	Type: 5	RData: FQDN alias
■ SOA	Type: 6	RData: Start Of Authority
■ PTR	Type: 12	RData: FQDN reverse lookup
■ MX	Type: 15	RData: FQDN mail exchange server
■ TXT	Type: 16	RData: Text
<ul><li>AAAA</li></ul>	Type: 28	RData: IPv6 address
<ul><li>DNSKEY</li></ul>	Type: 48	RData: PublicKey

RData: Improved access to https resources

#### **DNS RR A**

#### Resolves a hostname to an IPv4-address

\$ nslookup

> set query=**A** 

> www.fh-dortmund.de

Server: 192.168.1.1

Address: 192.168.1.1#53

Non-authoritative answer:

Name: www.fh-dortmund.de

Address: 193.25.16.26

#### **DNS RR AAAA**

#### Resolves a hostname to an IPv6-address

\$ nslookup

> set query=**AAAA** 

> www.google.de

Server: fe80::d487:62:9df4:e3dd%15

Address: fe80::d487:62:9df4:e3dd%15#53

Non-authoritative answer:

www.google.de has AAAA address 2a00:1450:4001:829::2003

#### **DNS RR CNAME**

Maps one domain name (an alias) to another(the canonical name).

```
Answers
    www.wdr.de: type CNAME, class IN, cname www.wdr.de.v1.edgekey.net
    Name: www.wdr.de
    Type: CNAME (Canonical NAME for an alias) (5)
    Class: IN (0x0001)
    Time to live: 44 (44 seconds)
    Data length: 27
    CNAME: www.wdr.de.v1.edgekey.net
    www.wdr.de.v1.edgekey.net: type CNAME, class IN, cname e8381.e6.akamaiedge.net
    Name: www.wdr.de.v1.edgekey.net
    Type: CNAME (Canonical NAME for an alias) (5)
    Class: IN (0x0001)
    Time to live: 20547 (5 hours, 42 minutes, 27 seconds)
    Data length: 22
    CNAME: e8381.e6.akamaiedge.net
```

#### **DNS RR MX**

Resolves the IP-address of the responsible mail server for this domain

\$ nslookup

> set query=**MX** 

> www.fh-dortmund.de

Server: 192.168.1.1

Address: 192.168.1.1#53

Non-authoritative answer:

fh-dortmund.de mail exchanger = 0 fhdortmund-de0e.mail.protection.outlook.com.

#### **DNS RR NS**

Resolves the authoritative nameserver for this domain

\$ nslookup

> set query=**NS** 

> fh-dortmund.de

Server: 192.168.1.1

Address: 192.168.1.1#53

Non-authoritative answer:

fh-dortmund.de nameserver = ns2.inwx.de.

fh-dortmund.de nameserver = ns3.inwx.eu.

fh-dortmund.de nameserver = ns.inwx.de.

#### **DNS RR SOA**

Provides administrative information about the zone

\$ nslookup
> set query=soa

> fh-dortmund.de

Server: 192.168.1.1

Address: 192.168.1.1#53

Non-authoritative answer:

#### fh-dortmund.de

origin = ns.inwx.de

mail addr = hostmaster.inwx.de

serial = 2023060500

refresh = 1800

retry = 900

expire = 21600

minimum = 600

## **DNS Domains & Zones**

 Domain the totality of descendant nodes of a certain node (a sub-tree)

#### Zone

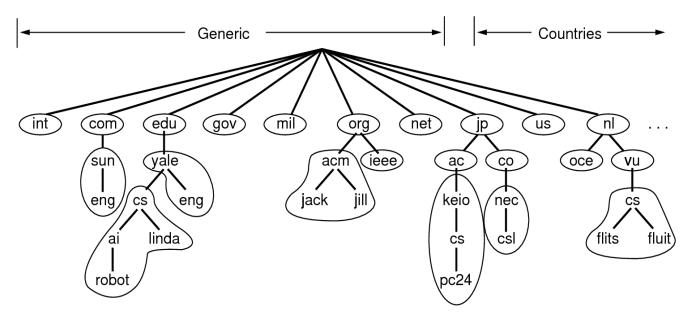
a portion of a domain, administered by an administrative entity. The boundaries of the zones are set by how authority is partitioned and distributed over the distributed database.

For administration, the (single) DNS tree is divided into zones. The administrative entity that has authority over a portion (zone) of the namespace can add, remove, or change labels within that zone. A zone begins at a certain node, and includes all nodes descended from it, except nodes belonging to sub-zones whose authority has previously been delegated to other entities.

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## **DNS Zones**

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The illustrations show examples of dividing individual domains into zones.

## Division of a domain in zones

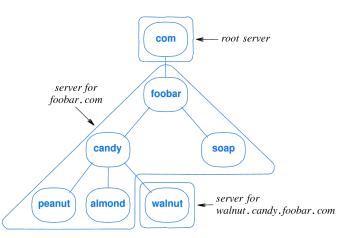
server for candy foobar com

candy foobar com

candy lamond walnut

A single server cannot manage all entries in a domain

- Dividing the domain into zones
- The division into zones is the responsibility of the driver of the domain.
- In the upper picture, a DNS Server for foobar.com and another DNS server for candy.foobar.com are used.
- In the lower picture, a DNS server for foobar.com and another DNS Server for walnut.candy.foobar.com are used.
- Structuring the domain into zones is not visible from the outside.



## **DNS zone file and FQDN**

- Stores all information related to the specified zone
  - Mapping between domain names and ip-addresses
  - Can store further information (DNSKEY, HTTPS)
- Text file format (defined in RFC 1034 and RFC 1035)
- Contains of line with either directives or resource records.
  - Directives are control entries that affect the rest of the zone file
  - o If entries end with ".", they are fully qualified, otherwise related to origin

#### FQDN

- o Fully qualified domain name
- o Specifies the **exact** location in the hierarchy tree
- Topmost layer is the root zone (.)

#### **DNS** zone file

\$ORIGIN example.com. ; designates the start of this zone file in the namespace \$TTL 3600 ; default expiration time (in seconds) IN SOA ns.example.com. username.example.com. ( 2020091025 7200 3600 1209600 3600 ) example.com IN NS ns ; ns.example.com is a nameserver for example.com example.com. example.com. IN NS ns.somewhere.example.; ns.somewhere.example is a backup nameserver for example.com IN MX 10 mail.example.com ; mail.example.com is the mailserver for example.com example.com. IN MX 20 mail2.example.com.; equivalent to above line, "@" represents zone origin example.com. IN A 192.0.2.1 ; IPv4 address for example.com IN AAAA 2001:db8:10::1 ; IPv6 address for example.com IN A 192.0.2.2 ; IPv4 address for ns.example.com ns IN AAAA 2001:db8:10::2 ; IPv6 address for ns.example.com IN CNAME example.com. ; www.example.com is an alias for example.com WWW wwwtest IN CNAME www ; www.example.com is another alias for www.example.com mail IN A 192.0.2.3 ; IPv4 address for mail.example.com IN A 192.0.2.4 ; IPv4 address for mail2.example.com mail2 IN AAAA 2001:db8:10::3; IPv6 address for file.example.com file

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#### **DNS** zone file

SORIGIN example.com. ; designates the start of this zone file in the namespace

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#### Fully qualified

```
$TTL 3600 ; default expiration time (in seconds)
example.com
                IN SOA ns.example.com. username.example.com. ( 2020091025 7200 3600 1209600 3600 )
example.com.
              ✓IN NS ns ; ns.example.com is a nameserver for example.com
example.com.
                IN NS ns.somewhere.example.; ns.somewhere.example is a backup nameserver for example.com
                IN MX 10 mail.example.com ; mail.example.com is the mailserver for example.com
example.com.
                IN MX 20 mail2.example.com.; equivalent to above line, "@" represents zone origin
example.com.
                IN A 192.0.2.1 ; IPv4 address for example.com
                IN AAAA 2001:db8:10::1 ; IPv6 address for example.com
                IN A 192.0.2.2 ; IPv4 address for ns.example.com
ns
                IN AAAA 2001:db8:10::2 ; IPv6 address for ns.example.com
                IN CNAME example.com. ; www.example.com is an alias for example.com
WWW
wwwtest
                IN CNAME www ; wwwtest.example.com is another alias for www.example.com
mail
                IN A 192.0.2.3 ; IPv4 address for mail.example.com
                IN A 192.0.2.4 ; IPv4 address for mail2.example.com
mail2
                IN AAAA 2001:db8:10::3; IPv6 address for file.example.com
file
```

#### **Root Server**

- Name server for the root zone of the DNS of the internet (.)
- Answers requests related to the TLD (returns a list of authoritative servers)
- 13 logical root servers in the internet
  - o a.root-server.net
  - b.root-server.net
  - 0 ...
  - o m.root-server.net



#### **Root Server II**

- One small zonefile
  - Contains names and IP addresses of authoritative DNS servers for each TLD
  - o Small (1.64MB as of 09 Jun 2023)
  - Changes infrequently (every couple of days)
- Q: How do I get to www.ub.fernuni-hagen.de?
- A: "I don't exactly know, but you should ask the .de TLD server at:
  - o a.nic.de, 194.0.0.53
  - s.de.net, 195.243.137.26
  - o and others

## **Root zone**

Online available at <a href="https://www.internic.net/domain/root.zone">https://www.internic.net/domain/root.zone</a></a> Entries for TLD .de

de.		172800	IN	NS	a.nic.de.
de.		172800	IN	NS	f.nic.de.
de.		172800	IN	NS	l.de.net.
de.		172800	IN	NS	n.de.net.
de.		172800	IN	NS	s.de.net.
de.		172800	IN	NS	z.nic.de.
<snip></snip>					
a.nic.de.	172800	IN	Α	194.0.0	).53
a.nic.de.	172800	IN	AAAA	2001:6	78:2:0:0:0:0:53
f.nic.de.	172800	IN	Α	81.91.1	L64.5
f.nic.de.	172800	IN	AAAA	2a02:5	68:0:2:0:0:0:53
<snip></snip>					

## **DNS** management

- The ICANN (Internet Corporation for assigned Names and Numbers) and the IANA (Internet Assigned Numbers Authority) manage these top level domains and assign subdomains to organizations, in Germany e.B. the DENIC (http://www.denic.de).
- The organization may then independently manage the labels below its subdomain. Either the next label is a computer name or the organization divides its domain into further subdomains.
- The FH Dortmund has the subdomain fh-dortmund.de. This could be further subdivided, e.B. into dvz.fh-dortmund.de, verw.fh-dortmund.de, informatik.fh-dortmund.de, etc. These subdomains are then managed by the institutions mentioned in the label.





The following table lists some of the top level domains:

Geographical Top-level domains		Top-level domains	
at	Austria	com	commercial companies
ch	Switzerland	edu	educational institutions
de	Germany	gov	government institutions
es	Spain	int	international organisations
fr	France	mil	military organisations
au	Australia	net	network operator
ca	Canada	org	non-commercial organisations

In 2000, ICANN decided to introduce seven new top level domains:

.aero aviation industry , .biz company , .coop cooperative organizations, .info without restriction , .museum museums , .name individuals , .pro lawyers, tax consultants, doctors

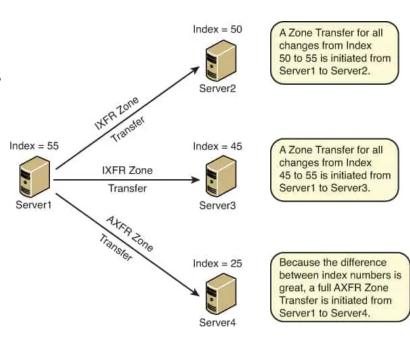
In the meantime, others have been added: e.B. . jobs, .mobi, .travel in 2005 In 2023, .zip was added (plus .dad, .phd, .prof, .esq, .foo, .mov, and .nexus.)

## **DNS** servers and zones

- The entity that has authority over a zone is responsible for maintaining the DNS servers corresponding to that zone (authoritative servers)
- There are 2 types of authoritative servers: primary (only one) and secondary.
- A primary server for one zone can be a slave for a different zone simultaneously
- The secondaries update their information automatically, consulting the primary (SOA registry) and doing a "zone transfer".
  - o RR type AXFR (Asynchronous Full Transfer Zone or Asynchronous Xfer Full Range
  - Zone transfer ensures replication of DNS DB
  - Transfer is done with Port 53/TCP (reliability is necessary)
  - The secondary server(s) updates the RRs from the primary server
  - Zone transfer allows distribution of the DNS
  - o Improved version is IXFR (Incremental zone transfer), which only transfers changed RRs

## **AFXR** and **IXFR** example

- All 4 servers are equal
  - Every server can response to requests
  - o Zone file contains (at least) 4 NS entries
- Changes are done on server 1 (master)
- Server 2-4 can be distributed all over the world
- If one server breaks, name resolution still possible



#### **Inserting records into DNS**

- Example: new startup "Network Utopia"
- register name networkuptopia.com at DNS registrar (e.g., Network Solutions)
  - o provide names, IP addresses of authoritative name server (primary and secondary)
  - o registrar inserts two RRs into com TLD server:

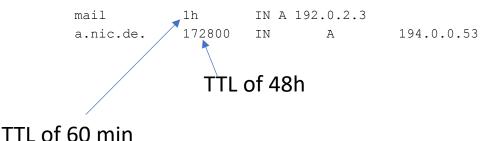
```
(networkutopia.com, dns1.networkutopia.com, NS)
(dns1.networkutopia.com, 212.212.212.1, A)
```

- Install DNS server and configure correct zone file
  - o create authoritative server type A record for www.networkuptopia.com
  - o type MX record for networkutopia.com
  - o Insert A, AAAA and further RR into the zone file
  - Keep entries up-to-date
  - o Increase serial no (in SOA RR) with each change

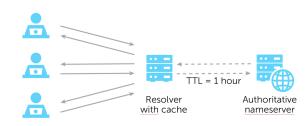


### **Recursive and caching only**

- On the client side, the resolution is usually mediated by a special server, called a recursive and caching-only nameserver, in which no information about any zones is configured.
- Caching-only server does entire name resolution process for the requesting client
- Server can use set TTL of a RR

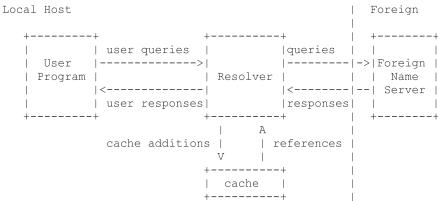


DNS cache time to live (TTL)



### Name resolving process

- In order for hosts to participate in the Domain Name System, they must use a Domain Name Resolver.
- This resolver software makes requests to the local domain name server to translate domain names into Internet addresses. Either the necessary information is available in the local name server or it must consult other name servers in order to provide the desired implementation.



#### **PTR**

- Used for Reverse DNS
  - Mapping of an ip-address to an hostname
- The domain in-addr.arpa (for IPv4) is used for PTR
  - o **ip6.arpa** used for IPv6
- Example:

lab.example.local has IP 192.168.2.42

- o 42.2.168.192.in-addr.arpa. 3600 IN PTR lab.example.local.

ipv6.example.local has IP 2001:db8::1

#### **Extensions**

- DNS is send in plain text over the channel
- Eavesdropping is possible
  - o Everyone with access to the communication is able to read the requests and responses
- "New" extensions prevent the plain text transmission
  - DoH, DNS over HTTPS
  - o DoT (RFC 7858), DNS over TLS
  - o DoQ (RFC 9250), DNS over QUIC

#### **DNS** and Wireshark

- Full support for DNS protocol
- Relevant filters:
  - o dns.a, dns.aaaa: Filter for dns responses
  - o dns.resp.name == FQDN: Filter for DNS responses with FQDN in it
  - dns.qry.name == FQDN: Filter for DNS requests for FQDN
  - o dns.cname: Filter CNAMEs

#### **Tools**

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```
dig
nslookup
host
$ host www.fh-dortmund.de
$ server:
```

www.fh-dortmund.de has address 193.25.16.26

```
└$ dig -t A www.fh-dortmund.de
          : <<>> DiG 9.10.6 <<>> -t A www.fh-dortmund.de
          ;; global options: +cmd
          ;; Got answer:
          ;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 17578
          ;; flags: qr rd ra; QUERY: 1, ANSWER: 1, AUTHORITY: 0, ADDITIONAL: 1
          :: OPT PSEUDOSECTION:
           EDNS: version: 0, flags:; udp: 512
          ;; QUESTION SECTION:
          :www.fh-dortmund.de.
                                       ΤN
                                               Α
         ;; ANSWER SECTION:
         www.fh-dortmund.de.
                                10
                                                      193,25,16,26
          ;; Query time: 74 msec
            SERVER: fe80::d487:62:9df4:e3dd%15#53(fe80::d487:62:9df4:e3dd%15)
            WHEN: Thu Jun 08 22:29:28 CEST 2023
          ;; MSG SIZE rcvd: 63
←$ nslookup www.fh-dortmund.de
                      fe80::d487:62:9df4:e3dd%15
Address:
                      fe80::d487:62:9df4:e3dd%15#53
Non-authoritative answer:
```

Address: 193.25.16.26

Name:

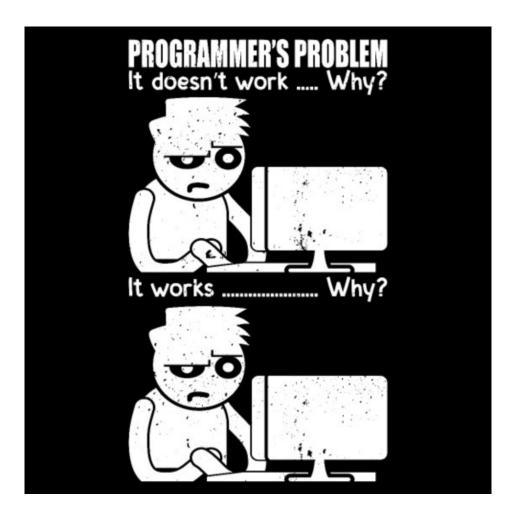
www.fh-dortmund.de

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# **Application Layer – HTTP and HTTPS**

Prof. Dr. Daniel Spiekermann Faculty of Computer Science

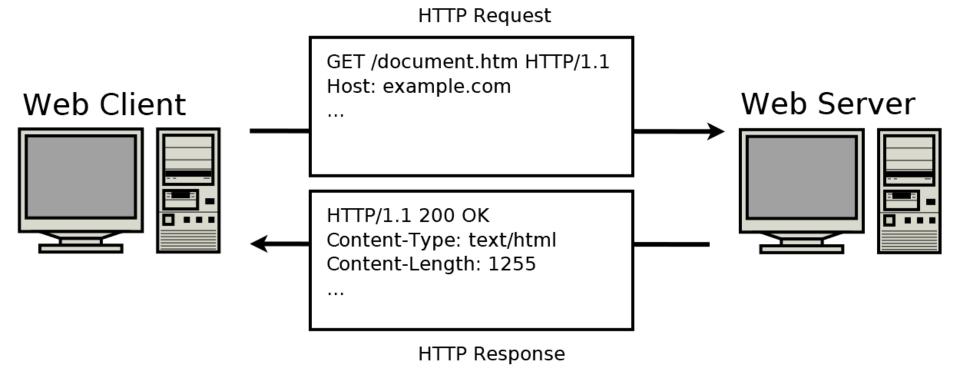
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- Hyper Text Transfer Protocol
- Developed by Tim Berners-Lee in 1989 at CERN
- Request/Response Protocol
- Stateless
- Uses Uniform Resource Locators (URL) for identifying and locating files
- Different versions
  - o HTTP/1.0 1996 obsoletes
  - o HTTP/1.1 1997
  - o HTTP/2 2015
  - o HTTP/3 2022
- Nowadays basement for various on-top protocols

#### **Client Server Communication**



#### **Communication**

- Client tries to connect to a server establishing a connection via TCP 3-Way HS
- Server listening on that port accepts the connection and then waits for a client's request message (Successful 3-way HS)
- Client sends its request to the server
- Upon receiving the request, the server sends back an HTTP response message (header plus a body if it is required).
  - The body of this message is typically the requested resource, although an error message or other information may also be returned.
- At any time client or server can close the connection (TCP FIN or RST)
- Closing a connection is usually advertised in advance by using one or more HTTP headers in the last request/response message sent to server or client.

#### Requests

#### HTTP defines a set of **request methods** to indicate the desired action

 GET: Requests a representation of the specified resource. Requests using GET should

only retrieve data.

HFAD: Asks for a response identical to a GET request, but without the response body

Submits an entity to the specified resource, often causing a change in state or side effects on the server. POST:

PUT: Replaces all current representations of the target resource with the request

payload.

DFIFTE: Deletes the specified resource.

CONNECT: Establishes a tunnel to the server identified by the target resource.

OPTIONS: Describes the communication options for the target resource.

TRACE: Performs a message loop-back test along the path to the target resource.

PATCH Applies partial modifications to a resource.

#### **Header fields**

- List of strings send from client or server
- Typically invisible to end-user
  - Only browser and server need the data
- Define various parameters
  - Encoding
  - User-agent
  - Caching and Aging
  - Compression
  - Content-statistics (length, type, hash,...)

### **Requests format**

- Requests contains
- Request line: Request method
  - o Requested URL
  - Protocol version
- At least 1 header field:
  - o Field name: Value

(GET /index.html HTTP1.1)

Connection: keep-alive

Cache-Control: max-age=0

Accept-Encoding: gzip, deflate, br

#### **Response format**

- Response message from server contains
- Status line:
  - Protocol version
  - Status code
  - Optional phrase
- At least 1 response header field:
  - o Field name: Value

HTTP/1.1 200 OK

```
Server: Apache/2.4.57 (Unix)
```

Last-Modified: Mon, 11 Jun 2007 18:53:14 GMT

Content-Type: text/html

### **Status Codes (RFC 9110)**

Status codes are issued by a server in response to a client's request

- 1xx: Informational response request was received process is continued
- 2xx: Successful
   Request was successfully received, understood, and accepted
- 3xx: Redirection further action needs to be taken in order to complete the request
- 4xx: Client error Request contains bad syntax or cannot be fulfilled
- 5xx: Server error
   Server failed to fulfil the request

#### **Wireshark and Status Codes**

FIOLOGOI	Length   Time
HTTP	481 HTTP/1.1 301 Moved Permanently
HTTP	1012 HTTP/1.1 404 Not Found (text/html)
HTTP	71 HTTP/1.1 404 Not Found (text/html)
HTTP	1012 HTTP/1.1 404 Not Found (text/html)
HTTP	985 HTTP/1.1 404 Not Found (text/html)
HTTP	379 HTTP/1.1 302 Found
HTTP	972 HTTP/1.1 201 Created (text/javascript)
HTTP	299 HTTP/1.1 408 Request Time-out (text/html)
HTTP	591 HTTP/1.1 301 Moved Permanently (text/html)

#### **Communication**

User:

10.0.0.1/index.html

#### **Communication**

#### **User:**

#### 10.0.0.1/index.html

#### **Browser:**

```
GET /index.html HTTP 1.1

Host: 10.0.0.1

User-Agent: Mozilla/5.0

Accept:
text/html,application/xhtml+xml,applicatio
n/xml;q=0.9,image/avif,image/webp,*/*;q=0.8

Accept-Language: en-GB,en;q=0.5 Accept-
Encoding: gzip, deflate, br Connection:
keep-alive
```

#### **Communication**

#### User:

10.0.0.1/index.html

#### **Browser:**

```
GET /index.html HTTP 1.1

Host: 10.0.0.1

User-Agent: Mozilla/5.0

Accept:
text/html,application/xhtml+xml,applicatio
n/xml;q=0.9,image/avif,image/webp,*/*;q=0.8

Accept-Language: en-GB,en;q=0.5 Accept-
Encoding: gzip, deflate, br Connection:
keep-alive
```

#### Server:

```
OK Date: Sat, 03 Jun 2023
14:14:23 GMT
Content-Type: text/html;
charset=UTF-8
Content-Length: 155
Last-Modified: Wed, 08 Jan
2003 23:11:55 GMT
Server: Apache/1.3.3.7
       (Red-Hat/Linux)
ETag: "3f80f-1b6-3e1cb03b"
Accept-Ranges: bytes
Connection: close
<html> <head> <title>An
Example Page</title>
</head> <body> It
works! </body> </html>
```

#### **Communication**

User:

.

10.0.0.1/index.html

Host: 10.0.0.1

**Browser:** 

User-Agent: Mozilla/5.0

GET /index.html HTTP 1.1

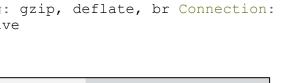
Accept:

text/html,application/xhtml+xml,applicatio n/xml;q=0.9,image/avif,image/webp,\*/\*;q=0.

8

Accept-Language: en-GB, en; q=0.5 Accept-Encoding: gzip, deflate, br Connection:

keep-alive



### Server:

**HTTP**/1.1 200 OK Date: Sat, 03 Jun 2023 14:14:23 GMT Content-Type: text/html; charset=UTF-8 Content-Length: 155 Last-Modifiéd: Wed, 08 Jan 2003 23:11:55 GMT Server: Apache/1.3.3.7 (Red-Hat/Linux) ETaq: "3f80f-1b6-3e1cb03b" Accept-Ranges: bytes Connection: close <html> <head> <title>An Example Page</title> </head> <body> It works! </body> </html>



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### Wireshark – Analyze HTTP

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Colorize Conversation
SCTP

Follow

Copy

Protocol Preferences
Decode As...

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TCP Stream
UDP Stream
DCCP Stream
TLS Stream
HTTP Stream

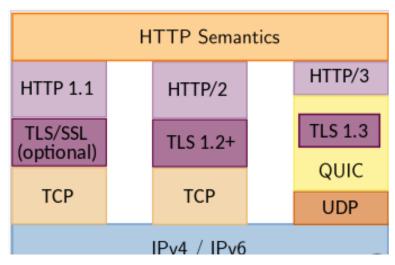
HTTD/2 Stroom

- Reconstructs the full communication process
  - Client request
  - Server response
- All header fields displayed

```
GET /index.html HTTP/1.1
Host: 127.0.0.1:8080
Connection: keep-alive
Cache-Control: max-age=0
sec-ch-ua: "Not.A/Brand"; v="8", "Chromium"; v="114", "Brave"; v="114"
sec-ch-ua-mobile: ?0
sec-ch-ua-platform: "macOS"
Upgrade-Insecure-Requests: 1
User-Agent: Mozilla/5.0 (Macintosh; Intel Mac OS X 10_15_7) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/114.0.0.0
Safari/537.36
Accept: text/html.application/xhtml+xml.application/xml;q=0.9,image/avif.image/webp,image/appq,*/*;q=0.8
Sec-GPC: 1
Accept-Language: de-DE.de
Sec-Fetch-Site: none
Sec-Fetch-Mode: navigate
Sec-Fetch-User: ?1
Sec-Fetch-Dest: document
Accept-Encoding: gzip, deflate, br
If-None-Match: "2d-432a5e4a73a80"
If-Modified-Since: Mon, 11 Jun 2007 18:53:14 GMT
HTTP/1.1 304 Not Modified
Date: Sat. 03 Jun 2023 14:18:48 GMT
Server: Apache/2.4.57 (Unix)
Last-Modified: Mon, 11 Jun 2007 18:53:14 GMT
ETag: "2d-432a5e4a73a80"
Accept-Ranges: bytes
Keep-Alive: timeout=5, max=100
Connection: Keep-Alive
```

### **Improvement of HTTP**

- HTTP/2 improves HTTP/1.1
  - Provides one or more bidirectional streams (prevents HOLB, Head of Line Blocking)
  - Less latency
  - Compression of header
- Current version is 3 (2022) as a revision of HTTP/2 (2015)
  - Uses QUIC+UDP instead of TCP



Keep-Alive: timeout=5, max=100

Connection: Keep-Alive



- HTTP transfers data in clear text
- Everyone with access to the connection can read the transferred data
- Eavesdropping of the connection gives access to all information
- Encryption does not prevent eavesdropping, but the clear text access

```
GET /index.html HTTP/1.1
Host: 127.0.0.1:8080
Connection: keep-alive
Cache-Control: max-age=0
sec-ch-ua: "Not.A/Brand"; v="8", "Chromium"; v="114", "Brave"; v="114"
sec-ch-ua-mobile: ?0
sec-ch-ua-platform: "macOS"
User-Agent: Mozilla/5.0 (Macintosh; Intel Mac OS X 10_15_7) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/114.0.0.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,image/apng,*/*;q=0.8
Accept-Language: de-DE, de
Sec-Fetch-Site: none
Sec-Fetch-Mode: navigate
Sec-Fetch-User: ?1
Sec-Fetch-Dest: document
Accept-Encoding: gzip, deflate, br
If-None-Match: "2d-432a5e4a73a80"
If-Modified-Since: Mon, 11 Jun 2007 18:53:14 GMT
Date: Sat, 03 Jun 2023 14:18:48 GMT
Server: Apache/2.4.57 (Unix)
Last-Modified: Mon, 11 Jun 2007 18:53:14 GMT
ETaq: "2d-432a5e4a73a80"
Accept-Ranges: bytes
```



```
.....aJT..~..X...C...,.>...?.gG.:8.d
....x0......<qo.-..1...:.***......+...0./...
     ....JJ......www.heise.de.....
JJ.....h2.http/1.1.....
.D..T..z...Q......xQ.......<qo.-..1...:.....3.$...zl....k....o
9..@.v`.....U....P...@.....FElwQ.....&.?r...r..T.W..Zv.-....Ij$..k...f)9.G.....i..{.9....p.8.....v....Y7.eb...
4r.^[.....+j.`...6...0n[..6..Q.....r.
.'I.*._.i.yY....BVp..j.W...!..1.......+/[.....E...:.s.IR.v...\.e-..%.*..]0..x......teP...
6P_...=j.o*.q.U...s..Y..[j..b.o.Tl.v...Z...L@}...Z...R.B^@':.9}.G....W...8..,0..mi,...P)X.ZJ.X.:5m...ns
3.#0{...}..<.8-.
Z...d(j...y{.I..2?.'.%......gQ.H..-...6...S..da.....B.J...
                                      ...D..&....0..i7).X'\.b...S.
4..'....I.....O../.Sw.#8D.+....G/8Po..k0.G......v..K..9..-Dl.e
```

### **HTTPs encryption**

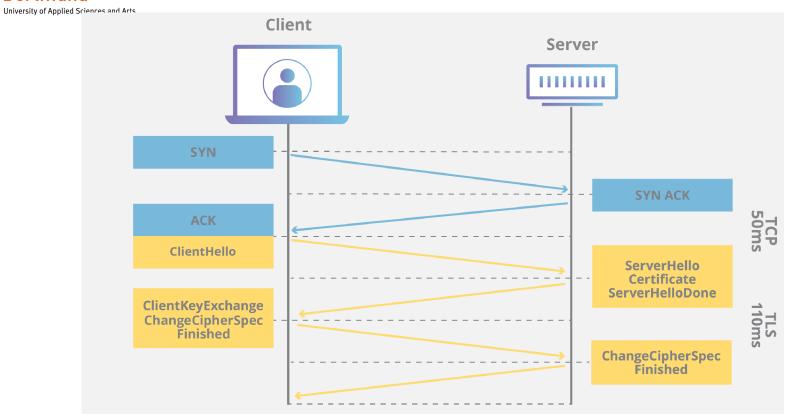
- HTTP over TLS
- HTTPS piggybacks HTTP entirely
- TLS (Transport Layer Security) after 3-way HS

	D 000111011011	1.10.0001	1 = 0.1.9.1.
192.168.179.203	193.99.144.85	TCP	78 59770 → 443 [SYN]
193.99.144.85	192.168.179.203	TCP	74 443 → 59770 [SYN,
192.168.179.203	193.99.144.85	TCP	66 59770 → 443 [ACK]
192.168.179.203	193.99.144.85	TLSv1.3	583 Client Hello
193.99.144.85	192.168.179.203	TCP	66 443 → 59770 [ACK]
193.99.144.85	192.168.179.203	TLSv1.3	1506 Server Hello, Cha
193.99.144.85	192.168.179.203	TCP	1506 443 → 59770 [ACK]
192.168.179.203	193.99.144.85	TCP	66 59770 → 443 [ACK]
192.168.179.203	193.99.144.85	TCP	66 59770 → 443 [ACK]
193.99.144.85	192.168.179.203	TLSv1.3	1506 Application Data
102 00 1// 05	102 160 170 202	TI Cv1 2	204 Application Data

#### **HTTPS** process

- Connection establishment from the client C to the server S over TCP 3-Way HS
- C sends cipher list (Client Hello)
- S authenticates with its certificate and chosen certificate (Server Hello, Certificate)
- C checks the cert and might send its own cert (Client Key Exchange, Client Certficate)
- C sends a random number (encrypted with pubkey of the server) or C and S use
   DH (Change Cipher Spec)
- All messages are encrypted with the calculated symmetric key

#### **Communication**



#### **Client Hello**

The message includes the client's TLS version, a random value, and a list of supported cipher suites.

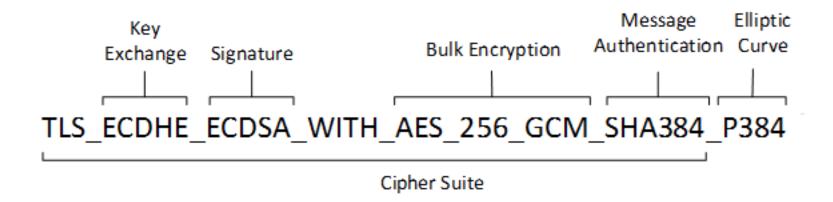
```
TLSv1.3 Record Layer: Handshake Protocol: Client Hello
    Content Type: Handshake (22)
    Version: TLS 1.0 (0x0301)
    Length: 512
  Handshake Protocol: Client Hello
      Handshake Type: Client Hello (1)
      Length: 508
      Version: TLS 1.2 (0x0303)
      Random: 16e86faf3ea23de2200413913c679b9d22c2f14308cea2ee1cc1a74a4214c482
      Session ID Length: 32
      Session ID: 784d1aae9ea109c0e5b510791772f986e77ca521b34b68271a96a0c4828818
      Cipher Suites Length: 42
    Cipher Suites (21 suites)
      Compression Methods Length: 1
    Compression Methods (1 method)
      Extensions Length: 393
    > Extension: Reserved (GREASE) (len=0)
    > Extension: server name (len=23)
    > Extension: extended_master_secret (len=0)
    > Extension: renegotiation info (len=1)
```

#### **Server Hello**

■ The server will respond by sending a Server Hello message which includes the TLS version, random value, and chosen cipher suite.

```
Transport Layer Security
  TLSv1.3 Record Layer: Handshake Protocol: Server Hello
     Content Type: Handshake (22)
     Version: TLS 1.2 (0x0303)
     Length: 122
    Handshake Protocol: Server Hello
        Handshake Type: Server Hello (2)
        Length: 118
        Version: TLS 1.2 (0x0303)
        Random: 78896013fa3021086e15296018f1d0d381c76220cd44a707c13e0502e2a3e819
        Session ID Length: 32
        Session ID: 784d1aae9ea109c0e5b510791772f986e77ca521b34b68271a96a0c4828818
        Cipher Suite: TLS AES 128 GCM SHA256 (0x1301)
        Compression Method: null (0)
        Extensions Length: 46
      > Extension: supported_versions (len=2)
      > Extension: key_share (len=36)
        [JA3S Fullstring: 771,4865,43-51]
```

### **Cipher Suite**



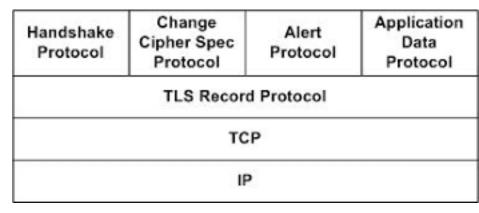
### **Client Key Exchange**

- After the client receives the server certificate, it will perform a series of verification steps to validate the certificate.
- If the result is positive, then the client will continue to the Client Key exchange.
- Per master Secret is encrpyted with certificate of the server
- Session secret is calculated on both sides with PMS, client and server random
- Both sides have now the same session secret (sym. enc)
- Client sends ClientHelloDone, Server sends ServerHelloDone



### **TLS Record Layer**

- Securing the TLS session
- Data is processed
  - Fragmentation (blocks of max 2<sup>14</sup> Byte)
  - Compression (no default)
  - MAC (Message Authentication Code) for integrity
  - Encryption
- Format of TLS Record:
- Content Type | Protokollversion Major (1 Byte) | Protokollversion Minor (1 Byte)
   Length (1 Short or 2 Byte)



### **HTTPS Content Types**

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- Content Types
- 20 Change Cipher Spec
- 21 Alert
- 22 Handshake
  - O Handshake Type:
    - 1: Client Hello
    - 2: Server Hello
    - 4: New Session Ticket
    - 11: Certificate
    - 12: Serer Key Exchange
    - 14: Server Hello Done
    - 16: Client Key Exchange
    - 22: Certificate Status
- 23 Application Data

```
TLSv1.2 Record Layer: Handshake Protocol: Certificate
    Content Type: Handshake (22)
    Version: TLS 1.2 (0x0303)
    Length: 2961
  Handshake Protocol: Certificate
      Handshake Type: Certificate (11)
      Length: 2957
      Certificates Length: 2954
    > Certificates (2954 bytes)
Transport Layer Security
TLSv1.2 Record Layer: Handshake Protocol: Certificate Status
    Content Type: Handshake (22)
    Version: TLS 1.2 (0x0303)
    Length: 479
  Handshake Protocol: Certificate Status
      Handshake Type: Certificate Status (22)
      Length: 475
      Certificate Status Type: OCSP (1)
      OCSP Response Length: 471
    > OCSP Response
  TLSv1.2 Record Layer: Handshake Protocol: Server Key Exchange
    Content Type: Handshake (22)
    Version: TLS 1.2 (0x0303)
    Length: 333
  Handshake Protocol: Server Key Exchange
      Handshake Type: Server Key Exchange (12)
      Length: 329
    > EC Diffie-Hellman Server Params
```

#### **TLS Layer**

- Handshake Protocol (Negotiation of session parameters)
- Alert Protocol (Management of alerts, warning or fatal. Fatal breaks connection)
- Change Cipher Spec Protocol (Changing to the negotiated cipher)
- Application Data Protocol (The transfer of the data over the TLS RR)

Handshake Protocol	Change Cipher Spec Protocol	Alert Protocol	Application Data Protocol
	TLS Record	d Protocol	
	тс	Р	
	IP	)	

#### **HTTPS and OSI**

- Precise classification difficult
  - HTTPS is not a new protocol, it is HTTP with TLS
  - Layer 6 translates or formats data for the application layer based on the semantics or syntax the application accepts. This layer also handles the **encryption** and decryption that the application layer requires.
  - o But TLS established a new session, which is part on L5

So TLS is sometimes assigned to L6, sometimes to L5, but sometimes to L7

OSI is not straightforward defined, it is a reference

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### **Additional slides**

Prof. Dr. Daniel Spiekermann Faculty of Computer Science